

Why Aren't More Users More Happy With Our VMs?



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Tratt

Warmup work in collaboration with:
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Software Development Team
2016-10-31

What to expect from this talk

The gap

What we tell users to expect

The gap

What we tell users to expect

The gap

What users experience

The gap

The gap
is bigger than we think

A stropky user? Or rightfully disappointed?

*"You told me I'd get
a 10x speed-up,
but I only saw 1.2x"*

Background

Dynamic ('JIT') compilation utilises
information not known statically

Is this just about 'dynamic typing'?

Is this statically or dynamically typed?

```
fn f(a: Option<i64>) {  
    match a {  
        Some(i) => ...,  
        None    => ...  
    };  
}
```

Is this just about 'dynamic typing'?

Is this statically or dynamically typed? Both!

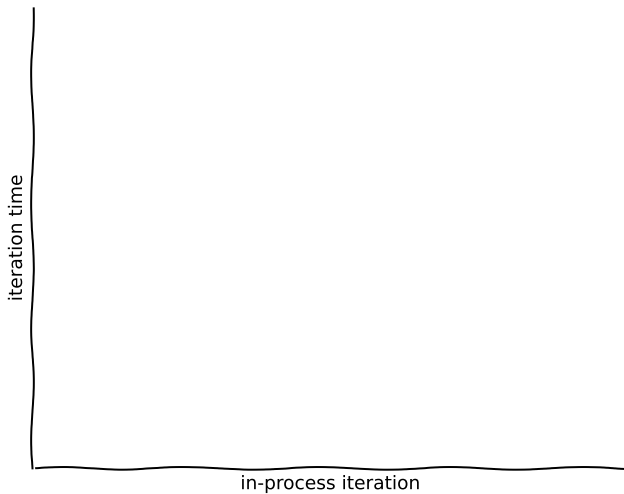
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fn f(a: Option<i64>) {  
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```

The best VMs are close in performance to,
and sometimes faster than,
gcc -O2

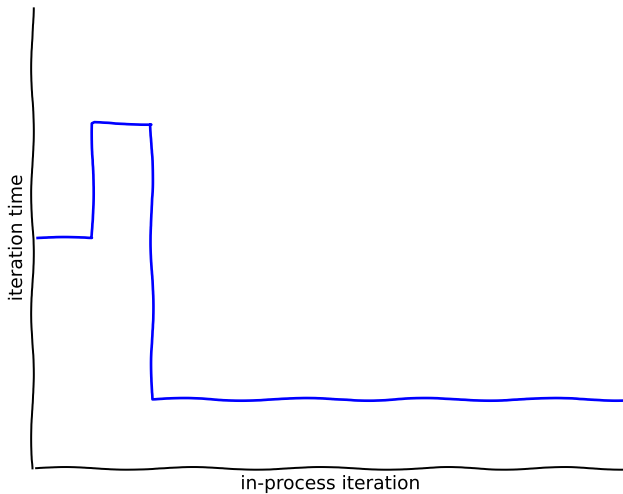
The best VMs are close in performance to,
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`gcc -O2`

What's being measured?

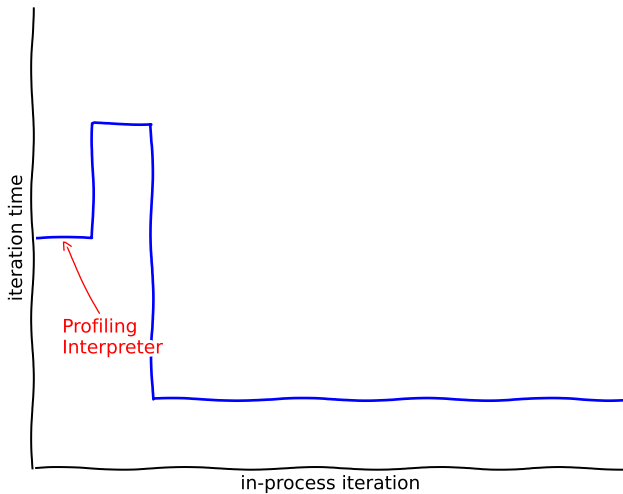
What our claims pertain to



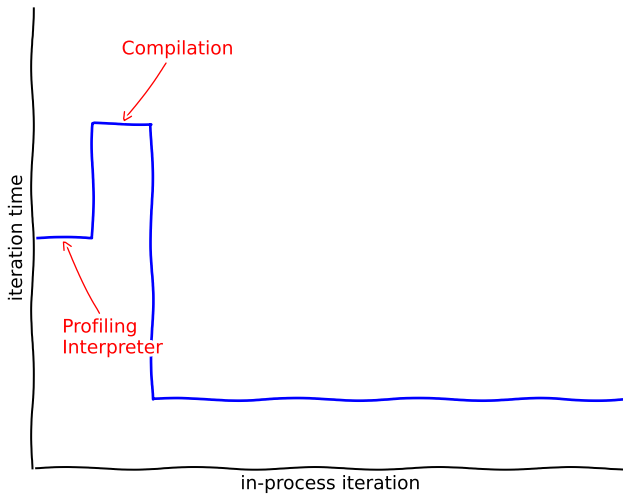
What our claims pertain to



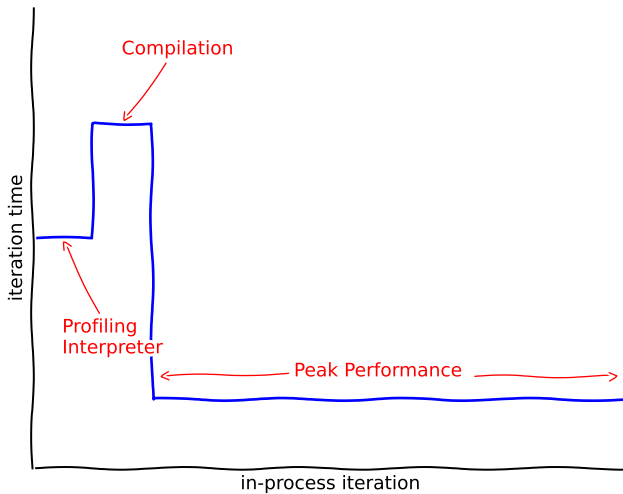
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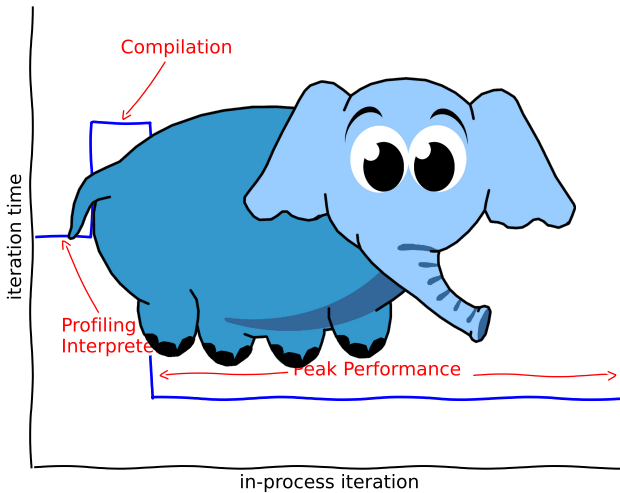
What our claims pertain to



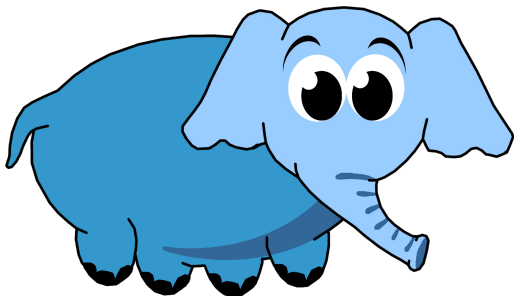
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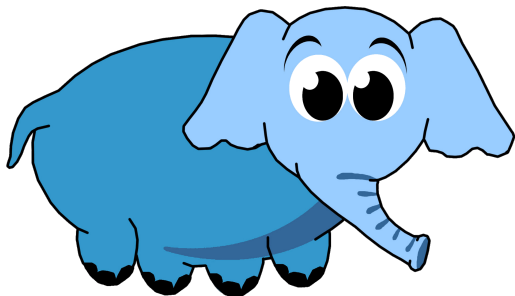


What our claims pertain to



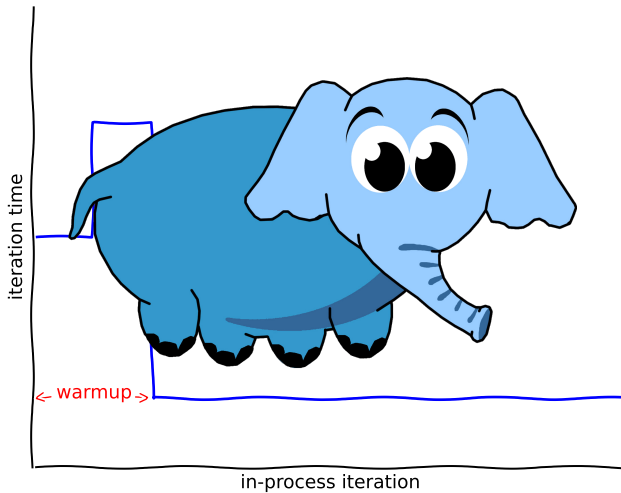
This is Barry

What our claims pertain to



This is Barry: the benchmarking elephant in the room

What our claims pertain to



Users *always* perceive warmup

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Maybe we should know how long it is?

The Warmup Experiment

Measure warmup of modern language
implementations

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Measure warmup of modern language
implementations

Hypothesis: Small, deterministic programs exhibit
classical warmup behaviour

Method 1: Which benchmarks?

The language benchmark games are perfect for us
(unusually)

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We removed any CFG non-determinism

We added checksums to all benchmarks

Method 2: How long to run?

2000 in-process iterations

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2000 in-process iterations

10 process executions

Method 3: VMs

- Graal-0.13
- HHVM-3.12.0
- JRuby/Truffle (git #f82ac771)
- Hotspot-8u72b15
- LuaJit-2.0.4
- PyPy-4.0.1
- V8-4.9.385.21
- GCC-4.9.3

Note: same GCC (4.9.3) used for all compilation

Method 4: Machines

- Linux-Debian8/i4790K, 24GiB RAM
- Linux-Debian8/i4790, 32GiB RAM
- OpenBSD-5.8/i4790, 32GiB RAM

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- Linux-Debian8/i4790K, 24GiB RAM
 - Linux-Debian8/i4790, 32GiB RAM
 - OpenBSD-5.8/i4790, 32GiB RAM
-
- Turbo boost and hyper-threading disabled
 - SSH blocked from non-local machines
 - Daemons disabled (cron, smtpd)

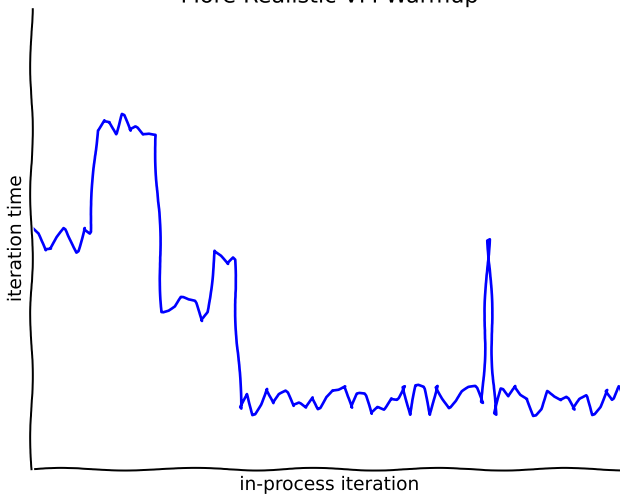
Benchmark runner: tries to control as many confounding variables as possible

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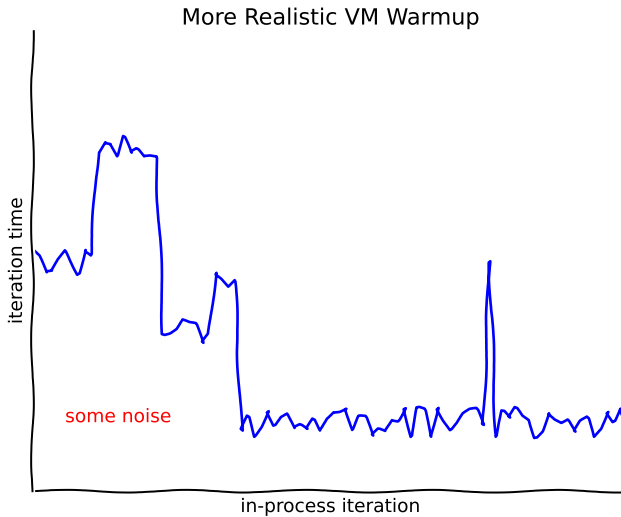
- Minimises I/O
- Sets fixed heap and stack ulimits
- Drops privileges to a 'clean' user account
- Automatically reboots the system prior to each proc. exec
- Checks `dmesg` for changes after each proc. exec
- Checks system at (roughly) same temperature for proc. execs
- Enforces kernel settings (tickless mode, CPU governors, ...)

Method 6: Expectations

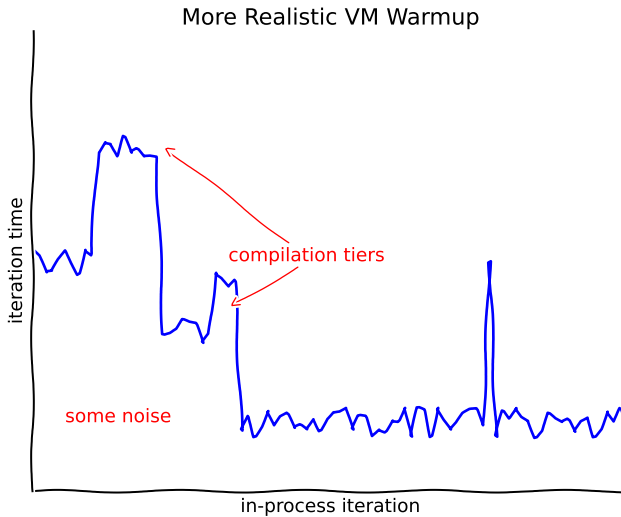
More Realistic VM Warmup



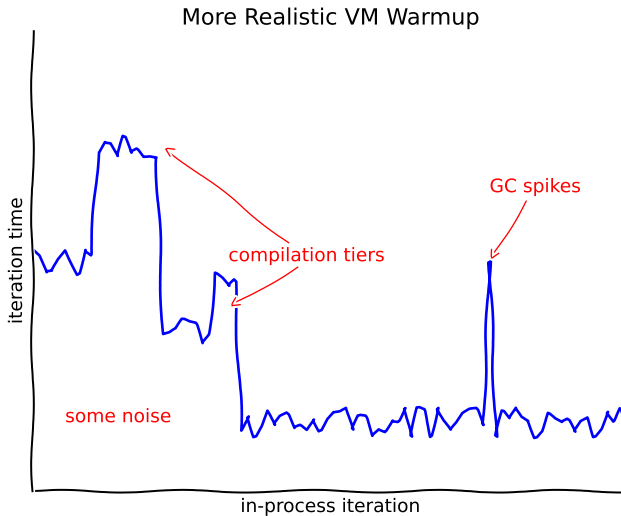
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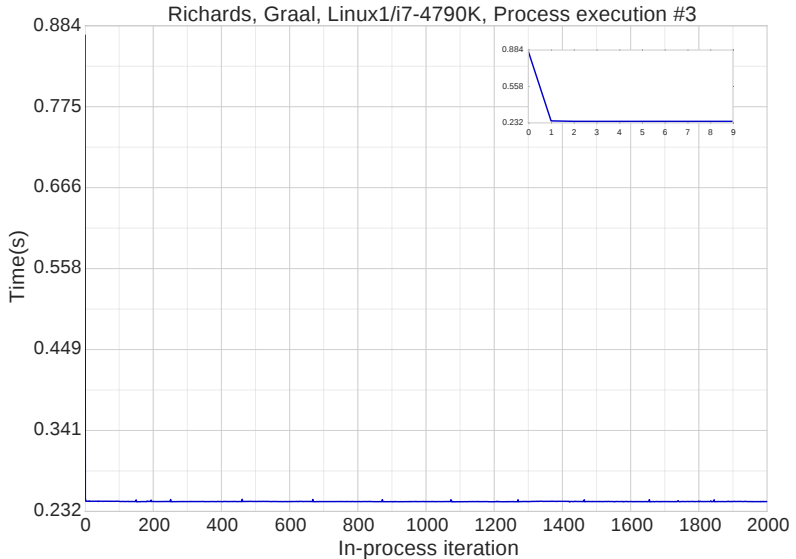


Method 6: Expectations

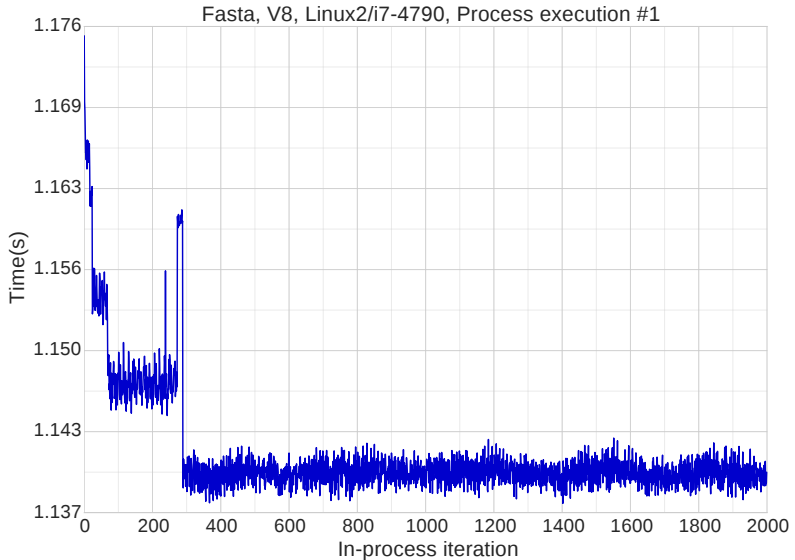


Preliminary results

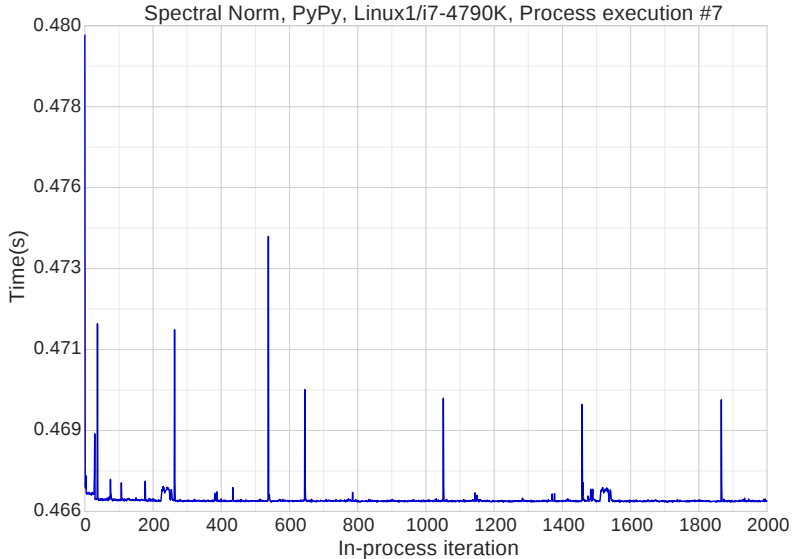
Classical Warmup



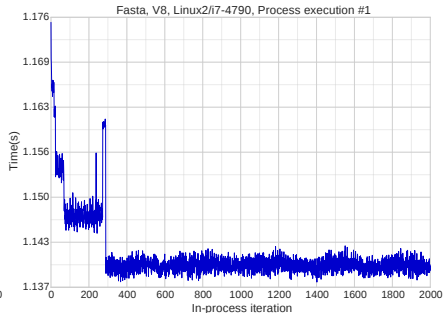
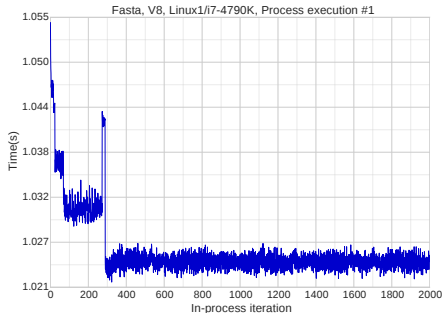
Classical Warmup



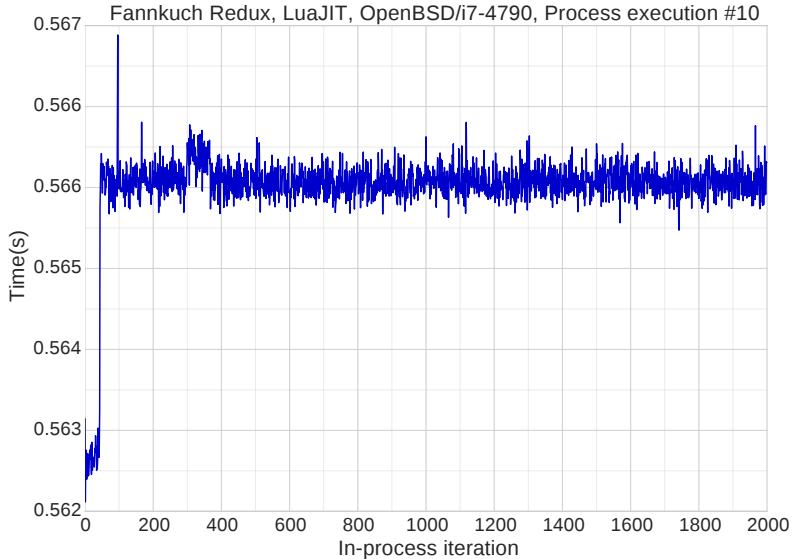
Classical Warmup



Classical Warmup across machines



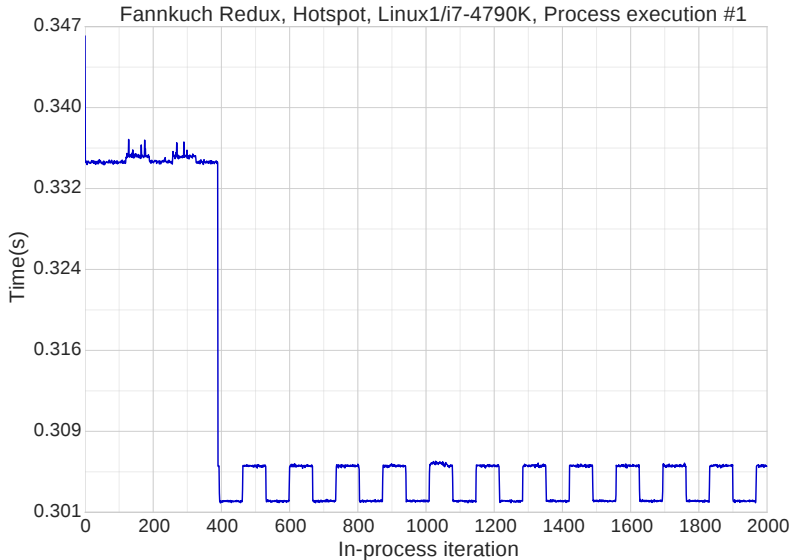
Slowdown



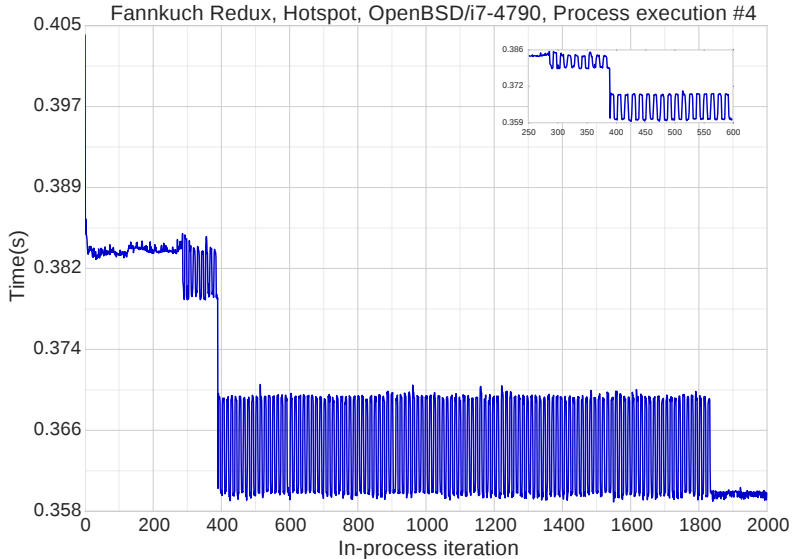
Slowdown



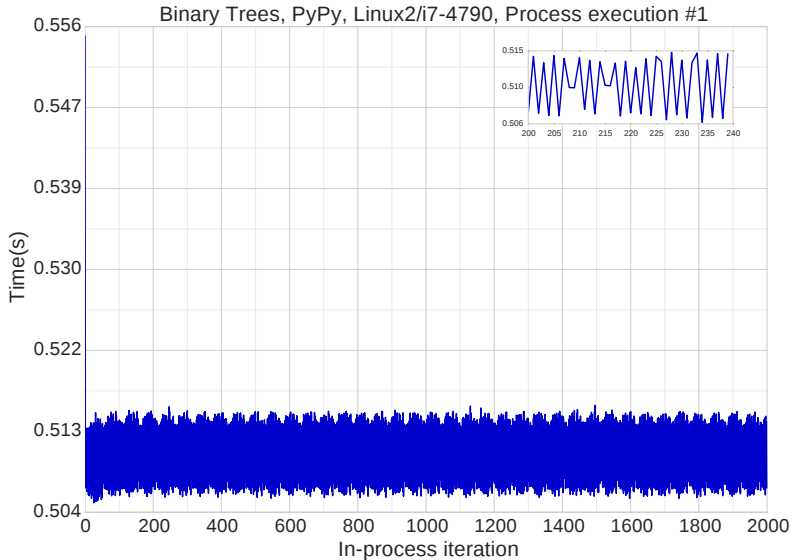
Cycles



Cycles



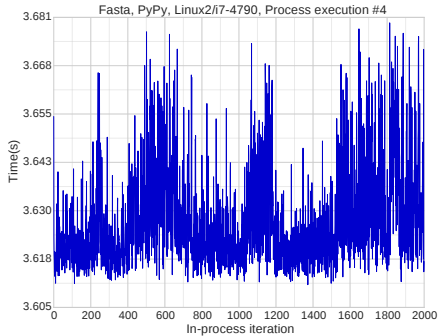
Cycles



Never-ending Phase Changes

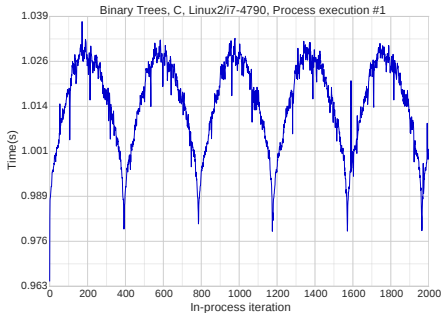


Inconsistent Process-executions



(Note: same machine)

Inconsistent Process-executions



(Note: different machines. Bouncing ball pattern
Linux-specific)

Classical warmup occurs for only:

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50% of process executions

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25% of (VM, benchmark) pairs

Classical warmup occurs for only:

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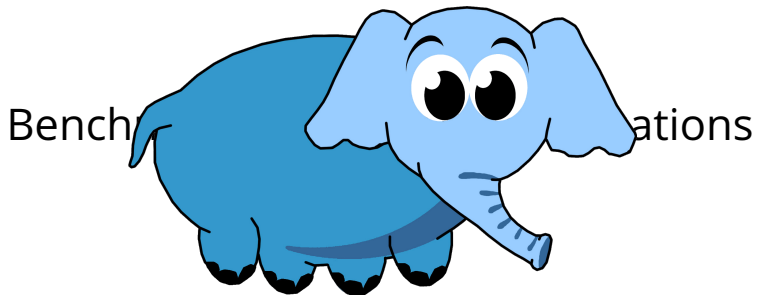
25% of (VM, benchmark) pairs

0% of benchmarks for all VMs, machines &
proc execs.

Benchmark suites

Benchmarks guide our optimisations

Benchmark suites



Benchmarks guide our optimisations

Are they complete guides?

A war story

Symptom: poor performance of a Pyston benchmark on PyPy

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Cause: RPython traces recursion

A war story

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Cause: RPython traces recursion

Fix: Check for recursion before tracing

A war story: the basis of a fix

```
diff --git a/rpython/jit/metainterp/pyjitpl.py b/rpython/jit/metainterp/pyjitpl.py
--- a/rpython/jit/metainterp/pyjitpl.py
+++ b/rpython/jit/metainterp/pyjitpl.py
@@ -951,9 +951,31 @@
     if warmrunnerstate.inlining:
         if warmrunnerstate.can_inline_callable(greenboxes):
+
+            # We've found a potentially inlinable function; now we need to
+            # see if it's already on the stack. In other words: are we about
+            # to enter recursion? If so, we don't want to inline the
+            # recursion, which would be equivalent to unrolling a while
+            # loop.
             portal_code = targetjitdriver_sd.mainjitcode
             return self.metainterp.perform_call(portal_code, allboxes,
                                                  greenkey=greenboxes)
+
+        inline = True
+        if self.metainterp.is_main_jitcode(portal_code):
+            for gk, _ in self.metainterp.portal_trace_positions:
+                if gk is None:
+                    continue
+                assert len(gk) == len(greenboxes)
+                i = 0
+                for i in range(len(gk)):
+                    if not gk[i].same_constant(greenboxes[i]):
+                        break
+            else:
+                # The greenkey of a trace position on the stack
+                # matches what we have, which means we're definitely
+                # about to recurse.
+                inline = False
+                break
+        if inline:
             return self.metainterp.perform_call(portal_code, allboxes,
                                                  greenkey=greenboxes)
```

A war story: mixed fortunes

Success: slow benchmark now 13.5x faster

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Failure: some PyPy benchmarks slow down

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Success: slow benchmark now 13.5x faster

Failure: some PyPy benchmarks slow down

Solution: allow *some* tracing into recursion

A war story: data

#unrollings	1	2	3	5	7	10	
hexiom2	1.3	1.4	1.1	1.0	1.0	1.0	
raytrace-simple	3.3	3.1	2.8	1.4	1.0	1.0	
spectral-norm	3.3	1.0	1.0	1.0	1.0	1.0	
sympy_str	1.5	1.0	1.0	1.0	1.0	1.0	
telco	4	2.5	2.0	1.0	1.0	1.0	
polymorphism	0.07	0.07	0.07	0.07	0.08	0.09	

<http://marc.info/?l=pypy-dev&m=141587744128967&w=2>

A war story: conclusion

The benchmark suite said 7 levels, so that's what I suggested

A war story: conclusion

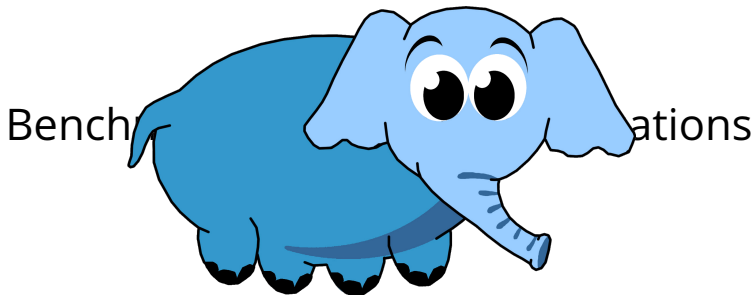
The benchmark suite said 7 levels, so that's what I suggested

Even though I doubted it was the right global value

Benchmark suites (2)

Benchmarks guide our optimisations

Benchmark suites (2)



Benchmarks guide our optimisations

Are they correct guides?

17 JavaScript benchmarks from V8

17 JavaScript benchmarks from V8

Let's make each benchmark run for 2000 iterations

Octane: pdf.js explodes

```
$ d8 run.js
Richards
DeltaBlue
Encrypt
Decrypt
RayTrace
Earley
Boyer
RegExp
Splay
NavierStokes
PdfJS
```

```
<--- Last few GCs --->
```

```
14907865 ms: Mark-sweep 1093.9 (1434.4) -> 1093.4 (1434.4) MB, 274.8 / 0.0 ms [allocation failure] [GC in old space]
14908140 ms: Mark-sweep 1093.4 (1434.4) -> 1093.3 (1434.4) MB, 274.4 / 0.0 ms [allocation failure] [GC in old space]
14908421 ms: Mark-sweep 1093.3 (1434.4) -> 1100.5 (1418.4) MB, 280.9 / 0.0 ms [last resort gc].
14908703 ms: Mark-sweep 1100.5 (1418.4) -> 1107.8 (1418.4) MB, 282.1 / 0.0 ms [last resort gc].
```

```
<--- JS stacktrace --->
```

```
==== JS stack trace =====
```

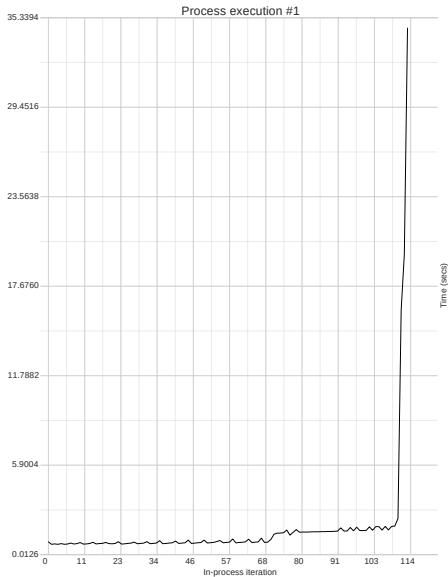
```
Security context: 0x20d333ad3ba9 <JS Object>
```

```
2: extractFontProgram(aka TypedParser_extractFontProgram) [pdfjs.js:17004] [pc=0x3a13b275421b] (this=0x3de358283)
3: new TypedFont [pdfjs.js:17216] [pc=0x3a13b2752078] (this=0x4603fbd9a9 <a TypedFont with map 0x1f822134f7e1>),
```

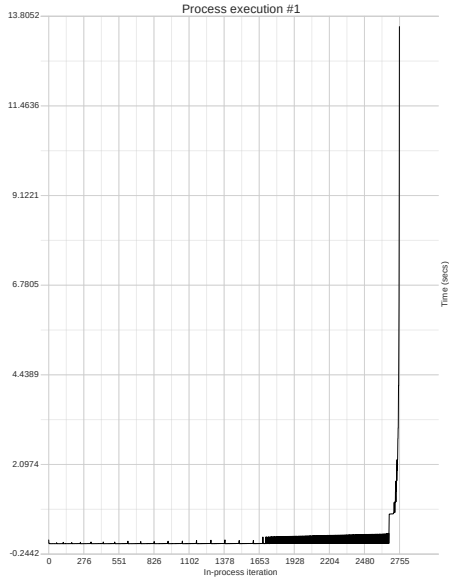
```
#
# Fatal error in CALL_AND_RETRY_LAST
# Allocation failed - process out of memory
#
```

```
zsh: illegal hardware instruction d8 run.js
```

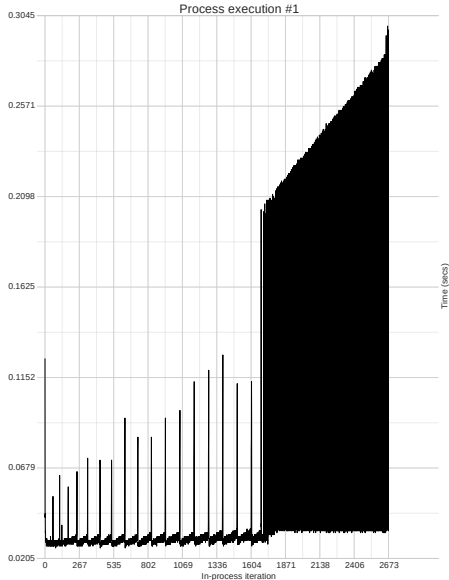
Octane: analysing pdf.js



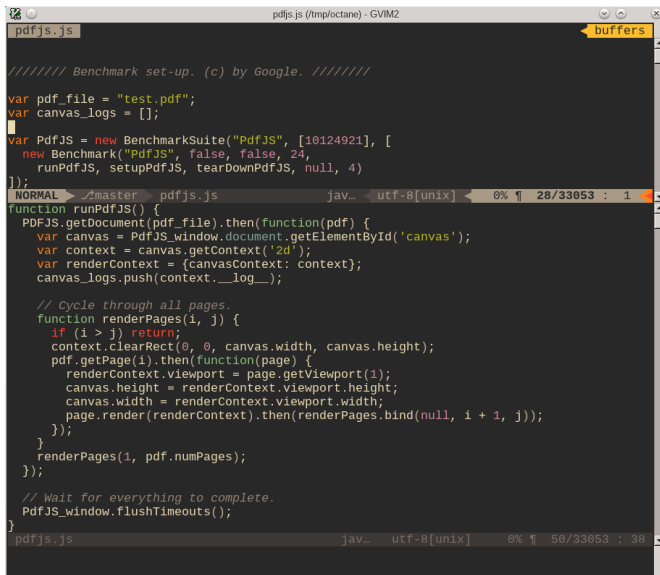
Octane: analysing pdf.js



Octane: analysing pdf.js



Octane: debugging



```
pdfjs.js

///////// Benchmark set-up. (c) by Google. //////////

var pdf_file = "test.pdf";
var canvas_logs = [];

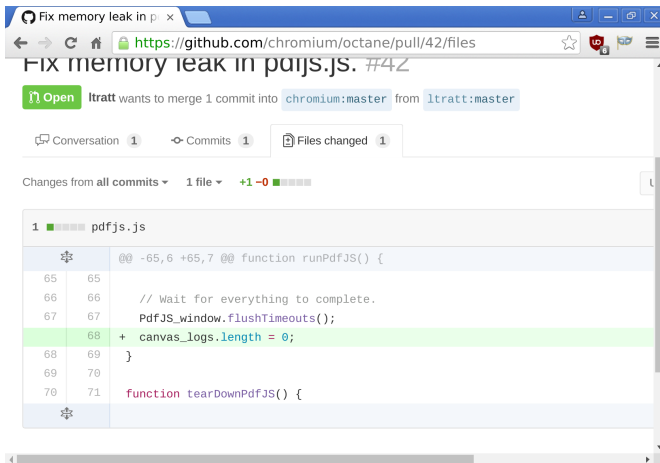
var PdfJS = new BenchmarkSuite("PdfJS", [10124921], [
  new Benchmark("PdfJS", false, false, 24,
    runPdfJS, setupPdfJS, tearDownPdfJS, null, 4)
]);

NORMAL 2 master pdfjs.js jav... utf-8[unix] 0% | 28/33053 : 1
function runPdfJS() {
  PDFJS.getDocument(pdf_file).then(function(pdf) {
    var canvas = PdfJS_window.document.getElementById('canvas');
    var context = canvas.getContext('2d');
    var renderContext = {canvasContext: context};
    canvas_logs.push(context.__log__);

    // Cycle through all pages.
    function renderPages(i, j) {
      if (i > j) return;
      context.clearRect(0, 0, canvas.width, canvas.height);
      pdf.getPage(i).then(function(page) {
        renderContext.viewport = page.getViewport(1);
        canvas.height = renderContext.viewport.height;
        canvas.width = renderContext.viewport.width;
        page.render(renderContext).then(renderPages.bind(null, i + 1, j));
      });
    }
    renderPages(1, pdf.numPages);
  });
}

// Wait for everything to complete.
PdfJS_window.flushTimeouts();
}
```

Octane: fixing



The screenshot shows a web browser window displaying a GitHub pull request. The address bar shows the URL `https://github.com/chromium/octane/pull/42/files`. The page title is "Fix memory leak in pdfjs.js. #42". Below the title, it says "ltratt wants to merge 1 commit into chromium:master from ltratt:master". There are tabs for "Conversation", "Commits", and "Files changed", with "Files changed" being the active tab. The changes section shows "Changes from all commits", "1 file", and a diff summary of "+1 -0". The diff for `pdfjs.js` is shown, with line 68 highlighted in green, indicating a new addition: `canvas_logs.length = 0;`. The surrounding code includes a `runPdfJS()` function and a `tearDownPdfJS()` function.

```
1 pdfjs.js
@@ -65,6 +65,7 @@ function runPdfJS() {
65 65
66 66     // Wait for everything to complete.
67 67     PdfJS_window.flushTimeouts();
68 + canvas_logs.length = 0;
68 69 }
69 70
70 71 function tearDownPdfJS() {
```

pdfjs isn't the only problem

pdfjs isn't the only problem

CodeLoadClosure also has a memory leak

pdfjs isn't the only problem

CodeLoadClosure also has a memory leak

zlib complains that Cannot enlarge memory
arrays in asm.js (a memory leak? I don't know)

pdfjs isn't the only problem

CodeLoadClosure also has a memory leak

zlib complains that Cannot enlarge memory arrays in `asm.js` (a memory leak? I don't know)

Timings are made with a non-monotonic
microsecond timer

Summary

Why aren't more users more happy with
our VMs?

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our VMs?

My thesis: our benchmarking *and* our
benchmarks have misled us

What we can do

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- 1 Run benchmarks for longer to uncover issues.

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- 2 Accept that peak performance may not occur.

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- 4 Stop over-training on small benchmark suites.

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- 5 The more benchmarks, the better.

What we can do

- 1 Run benchmarks for longer to uncover issues.
- 2 Accept that peak performance may not occur.
- 3 Always report warmup time.
- 4 Stop over-training on small benchmark suites.
- 5 The more benchmarks, the better.
- 6 Focus on predictable performance.

VM Warmup Blows Hot and Cold

E. Barrett, C. F. Bolz, R. Killick, V. Knight, S. Mount and L. Tratt.

Rigorous Benchmarking in Reasonable Time

T. Kalibera and R. Jones

Specialising Dynamic Techniques for Implementing the Ruby Programming Language

C. Seaton (Chapter 4)

Quantifying performance changes with effect size confidence intervals

T. Kalibera and R. Jones

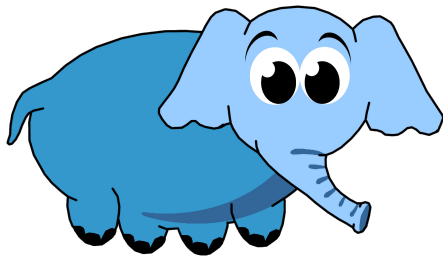
Thanks

- EPSRC: *COOLER* and *Lecture*.
- Oracle: various.

K Whiteford for Barry, the Benchmarking Elephant in the Room

Thanks for listening

Thanks for listening



And don't forget Barry